

CPSC 406 : Computational Optimization

CONVEX FUNCTIONS

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- A function $f: C \rightarrow \mathbb{R}$ is convex if $C \subseteq \mathbb{R}^n$ is convex and

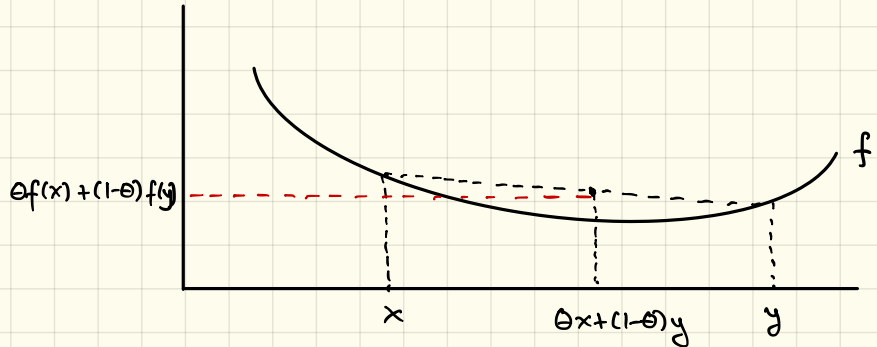
$$f(\theta x + (1-\theta)y) \leq \theta f(x) + (1-\theta)f(y)$$

for all $x, y \in C$ and $\theta \in [0, 1]$.

- f is strictly convex if for all $x, y \in C$ and $\theta \in (0, 1)$,

$$f(\theta x + (1-\theta)y) < \theta f(x) + (1-\theta)f(y)$$

- f is concave if $-f$ is convex



EXAMPLES

CONVEX FUNCTIONS

- affine : $a^T x + \beta$ for any $a \in \mathbb{R}^n$, $\beta \in \mathbb{R}$
- exponential: $e^{\alpha x}$ for any $\alpha \in \mathbb{R}$
- powers: x^α on \mathbb{R}_{++} for all $\alpha \geq 1$ or $\alpha \leq 0$
- abs val: $|x|^p$ for all $p \geq 1$
- neg entropy: $x \log x$ on \mathbb{R}_{++}
- norms (use triangle ineq. to show)

CONCAVE FUNCTIONS

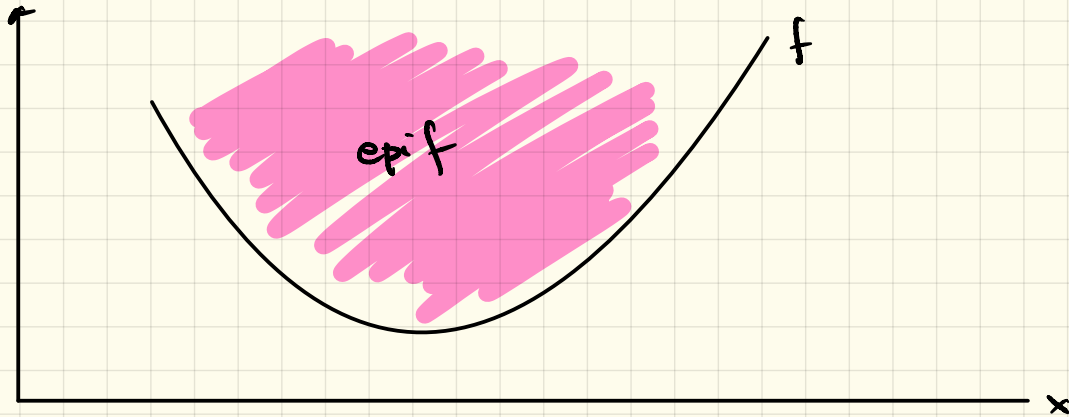
- affine : (see above)
- powers: x^α on \mathbb{R}_{++} for all $0 \leq \alpha \leq 1$
- logarithm: $\log x$ on \mathbb{R}_{++}

Equivalent Def'n for convexity of functions:

Def'n (II) A function $f: \mathbb{R}^n \rightarrow \mathbb{R}$ is convex if the epigraph,

$$\text{epi}f = \{ (x, \alpha) \in \mathbb{R}^n \times \mathbb{R} \mid f(x) \leq \alpha \} \subseteq \mathbb{R}^{n+1}$$

is a convex set.



RESTRICTION TO LINES

The function $f: C \rightarrow \mathbb{R}$ ($C \subseteq \mathbb{R}^n$ conv) is convex iff

$$\psi(\alpha) = f(x + \alpha d)$$

is convex for all $x \in C$ and $d \in \mathbb{R}^n$.

$$\psi(\alpha) = f(x + \alpha d)$$



[Recall discussion on line search]

OPERATIONS THAT PRESERVE CONVEXITY

- non-negative multiple

αf is convex if f is convex and $\alpha \geq 0$.

- sum (including infinite sums)

$f_1 + f_2$ is convex if f_1, f_2 are convex.

- composition with affine function

$f(Ax+b)$ is convex if f is convex

Examples

- log barrier: $f(x) = -\sum_{i=1}^m \log(a_i^T x - b_i)$ convex over $\{x \mid a_i^T x > b_i\} \forall i$

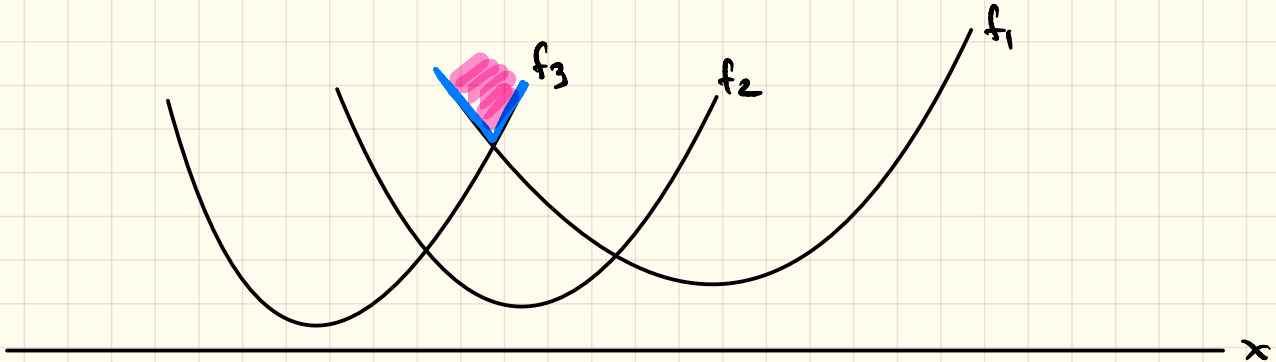
- norm of affine function: $f(x) = \|Ax - b\|$ [composition w/ affine + Δ -inequality]

- $f(x_1, x_2, x_3) = e^{x_1 - x_2 + x_3} + e^{2x_2} + x_1$

Suppose f_i , $i=1 \dots m$, are convex functions.

Then $f(x) = \max \{ f_1(x), f_2(x), \dots, f_m(x) \}$

is convex.



GLOBAL OPTIMALITY

$$(P) \underset{x}{\text{minimize}} f(x) \text{ subj to } x \in C$$

where f is a convex function and C is a convex set.

If x^* is a local minimizer of (P) $[f(x^*) \leq f(x) \forall x \in C \cap B_\epsilon(x^*)]$
it is also a global minimizer of (P) $[f(x^*) \leq f(x) \forall x \in C]$

Proof Suppose \bar{x} is a local, but not global, minimizer. Then,

• there exists $y \in C$ such that $f(y) < f(\bar{x})$.

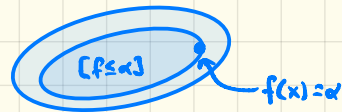
$$\begin{aligned} \bullet f(\theta \bar{x} + (1-\theta)y) &\leq \theta f(\bar{x}) + (1-\theta)f(y) \\ &< \theta f(\bar{x}) + (1-\theta)f(\bar{x}) \\ &= f(\bar{x}) \end{aligned}$$

which contradicts hypothesis.

LEVEL SETS

The level set of a function $f: C \rightarrow \mathbb{R}$ is the set

$$[f \leq \alpha] := \{x \in C \mid f(x) \leq \alpha\}$$



If f is convex \Rightarrow all of its level sets are convex.
(\uparrow note that not \Leftarrow)

Proof

Take $x, y \in [f \leq \alpha]$. Then $f(x) \leq \alpha$ and $f(y) \leq \alpha$.

Because f is convex, for $\theta \in [0, 1]$

$$f(\theta x + (1-\theta)y) \leq \theta f(x) + (1-\theta)f(y) \leq \theta \alpha + (1-\theta)\alpha \leq \alpha$$

Thus, $\theta x + (1-\theta)y \in [f \leq \alpha]$.

Corollary The set of minimizers of (P) is convex.

FIRST-ORDER CHARACTERIZATION

Let $f: C \rightarrow \mathbb{R}$ be cont. diff over $C \subseteq \mathbb{R}^n$ (cvx). Then f is convex iff

$$f(x) + \nabla f(x)^T(y-x) \leq f(y) \quad \text{for all } x, y \in C$$

Proof sketch:

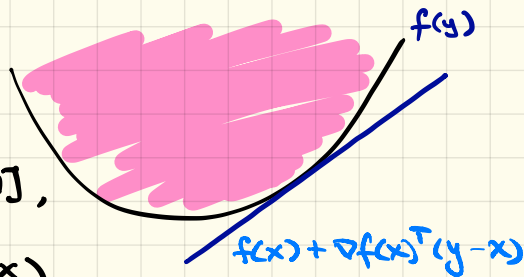
By convexity of f , $\forall x, y \in C$, $\theta \in [0, 1]$,

$$f(\theta y + (1-\theta)x) \leq \theta f(y) + (1-\theta)f(x)$$

$$\Rightarrow \frac{f(\theta y + (1-\theta)x) - f(x)}{\theta} \leq f(y) - f(x)$$

Take $\lim_{\theta \searrow 0}$, then

$$\underbrace{f'(x; y-x)}_{= \nabla f(x)^T(y-x)} \leq f(y) - f(x) \quad \left. \vphantom{\frac{f'(x; y-x)}{= \nabla f(x)^T(y-x)}} \right\} \Rightarrow f(y) \geq f(x) + \nabla f(x)^T(y-x)$$



FIRST-ORDER SUFFICIENCY

By convexity

$$f(z) \geq f(x) + \nabla f(x)^T (z-x) \quad \forall x, z.$$

Suppose $x = x^* \Rightarrow \nabla f(x^*) = 0$. Then

$$f(z) \geq f(x^*) \quad \forall z.$$

Second-Order Characterization

For $f: \mathbb{R}^n \rightarrow \mathbb{R}$ twice continuously differentiable on an open convex set $C \subseteq \mathbb{R}^n$, then f is convex over C iff

$$\nabla^2 f(x) \succeq 0 \quad \forall x \in C$$

Proof sketch (only \Leftarrow)

- By the second-order Taylor expansion of f :

$$f(y) = f(x) + \nabla f(x)^T (y-x) + \frac{1}{2} (y-x)^T \nabla^2 f(z) (y-x)$$

for $x, y \in C$ and some $z \in [x, y]$.

- Because $\nabla^2 f(x) \succeq 0$, $(y-x)^T \nabla^2 f(z) (y-x) \geq 0$

- Hence,

$$f(y) \geq f(x) + \nabla f(x)^T (y-x) \quad \forall x, y \in C$$

and f is convex over C .

Example $f(x) = x^\alpha$ for $x \geq 0$ $\alpha \geq 0$

- Differentiating twice over $x > 0$:

$$f''(x) = \alpha(\alpha-1)x^{\alpha-2}$$

- $f''(x) \geq 0$ if $\alpha \geq 1$
- $f''(x) \leq 0$ if $0 \leq \alpha \leq 1$