

CPSC 406: Computational Optimization

LINEAR CONSTRAINTS

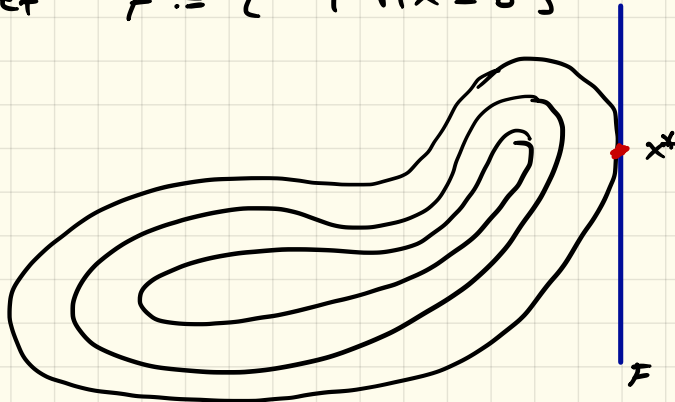
	February	13,	2018
revised	February	26,	2020
revised	March	2,	2023

LINEARLY - CONSTRAINED OPTIMIZATION

$$(P) \quad \underset{x \in \mathbb{R}^n}{\text{minimize}} \quad f(x) \quad \text{subj to} \quad Ax = b$$

- $f: \mathbb{R}^n \rightarrow \mathbb{R}$ continuously differentiable (once or twice, as needed)
- $A = \begin{matrix} \boxed{} \\ n \end{matrix} m$ with $\text{rank}(A) = m$
- minimum value finite and attained

Feasible set $F := \{x \mid Ax = b\}$



EXAMPLE (IN CLASS)

$$\text{minimize } \frac{1}{2}(x_1^2 + x_2^2) \quad \text{subj to } x_1 + x_2 = 1$$

OPTIMALITY CONDITIONS

Define "reduced" objective for any particular sol'n \bar{x} and basis Z :

$$f_z(p) := f(\bar{x} + Zp)$$

$$\nabla f_z(p) = Z^T \nabla f(\bar{x} + Zp) \quad [\text{reduced gradient}]$$

Let p^* be solution and set $x^* = \bar{x} + Zp^*$. Then p^* is optimal only if

$$\nabla f_z(p^*) = 0 \iff Z^T \nabla f(x^*) = 0 \iff \nabla f(x^*) \in \text{null}(Z^T)$$

Fundamental subspaces in \mathbb{R}^n associated with A and Z :

$$\text{range}(A^T) \oplus \text{null}(A) = \mathbb{R}^n$$

$$\text{null}(Z^T) \oplus \text{range}(Z) = \mathbb{R}^n$$

$$\text{Thus, } \nabla f(x^*) \in \text{null}(Z^T) \iff \nabla f(x^*) \in \text{range}(A^T)$$

$$\iff \exists y \text{ st } \nabla f(x^*) = A^T y$$

FIRST-ORDER NECESSARY CONDITIONS

A point x^* is a local minimizer of (P) only if

there exists an m -vector y such that

[optimality]

$$\nabla f(x^*) = A^T y \equiv \sum_{i=1}^m a_i y_i$$

$$A = \begin{bmatrix} a_1^T \\ a_2^T \\ \vdots \\ a_m^T \end{bmatrix}$$

[feasibility]

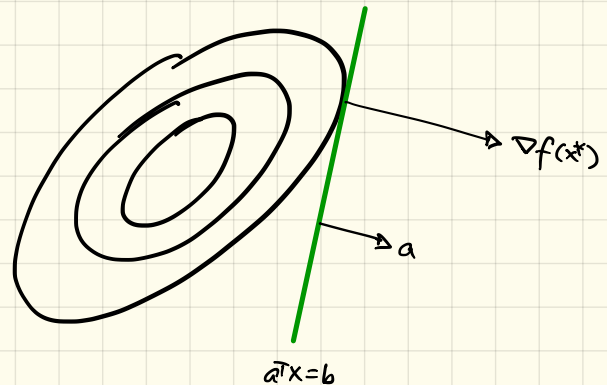
$$Ax^* = b$$

Equivalent to

$$z^T \nabla f(x^*) = 0 \iff \nabla f(x^*)^T p = 0 \quad \forall p \in \text{null}(A)$$

$$Ax^* = b$$

The vector y is sometimes referred to as "Lagrange multipliers".



SECOND-ORDER OPTIMALITY

$$f_z(p) := f(\bar{x} + Zp), \quad \nabla f_z(p) = Z^T \nabla f(\bar{x} + Zp), \quad \nabla^2 f_z(p) = Z^T \nabla^2 f(\bar{x} + Zp) Z$$

Necessary 2nd-order optimality: x^* is a local minimizer only if

$$\left. \begin{array}{l} Ax^* = b \\ Z^T \nabla f(x^*) = 0 \\ Z^T \nabla^2 f(x^*) Z \succ 0 \end{array} \right\} \begin{array}{c} \text{or} \\ \iff \end{array} \left\{ \begin{array}{l} Ax^* = b \\ \nabla f(x^*) = A^T y \text{ for some vector } y \\ p^T \nabla^2 f(x^*) p \geq 0 \quad \forall p \in \text{null}(A) \end{array} \right.$$

Sufficient 2nd-order optimality: x^* is a local minimizer if

$$\left. \begin{array}{l} Ax^* = b \\ Z^T \nabla f(x^*) = 0 \\ Z^T \nabla^2 f(x^*) Z \succ 0 \end{array} \right\} \begin{array}{c} \text{or} \\ \iff \end{array} \left\{ \begin{array}{l} Ax^* = b \\ \nabla f(x^*) = A^T y \text{ for some vector } y \\ p^T \nabla^2 f(x^*) p > 0 \quad \forall 0 \neq p \in \text{null}(A) \end{array} \right.$$

EXAMPLE : LEAST-NORM SOLUTIONS

minimize $\|x\|_2$ subj to $Ax=b$ (underdetermined)

Can take $f(x) = \frac{1}{2}\|x\|^2$. First-order optimality:

$$\left. \begin{array}{l} x = A^T y \text{ for some } y \\ Ax = b \end{array} \right\} \Leftrightarrow \begin{pmatrix} -I & A^T \\ A & 0 \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} 0 \\ b \end{pmatrix}$$

EXAMPLE Find a minimal-norm solution to

$$x_1 + x_2 + \dots + x_n = 1$$

$$\begin{pmatrix} -I & e \\ e^T & 0 \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$$\begin{array}{l} -x + ey = 0 \\ e^T x = 1 \end{array} \Rightarrow x = ey$$

$$\Rightarrow e^T ey = ny = 1 \quad y = \frac{1}{n} \Rightarrow \underline{x = \frac{1}{n}e}.$$

REDUCED GRADIENT METHOD

$$\begin{array}{ll} \text{minimize} & f(x) \\ \text{st} & Ax=b \end{array}$$

\Leftrightarrow

$$\text{minimize } f_z(\bar{x} + Zp)$$

x_0 feasible ; Z basis for $\text{null}(A)$

for $k=0, 1, 2, \dots$

- compute $g := \nabla f(x_k)$

- compute $H := \nabla^2 f(x_k)$

- solve $Z^T H Z p_k = -Z^T g$

ie $\nabla_z^2 f(x_k) p_k = -\nabla_z f(x_k)$

- linesearch on $f(x_k + \alpha Z p_k)$

- $x_{k+1} = x_k + \alpha_k Z p_k$

OBTAINING A NULL-SPACE BASIS

Permute variables (columns of A) so that

$$A = \begin{bmatrix} B & N \\ m & n-m \end{bmatrix} \quad \text{where } B \text{ nonsingular}$$

$B \equiv$ "Basic"
 $N \equiv$ "Non-basic"

Then feasibility requires

$$Ax = b \Leftrightarrow \begin{bmatrix} B & N \end{bmatrix} \begin{pmatrix} x_B \\ x_N \end{pmatrix} = b \Leftrightarrow Bx_B + Nx_N = b$$

Basic (x_B) and Non-Basic (x_N) variables:

- x_N free to move
- x_B uniquely determined by x_N , i.e., $x_B = B^{-1}(b - Nx_N)$

Constructing a null-space matrix:

$$Z = \begin{bmatrix} -B^{-1}N \\ I \end{bmatrix} \Rightarrow AZ = \begin{bmatrix} B & N \end{bmatrix} \begin{bmatrix} -B^{-1}N \\ I \end{bmatrix} = 0$$